

Irreducibility Testing of Nilpotent Matrix Groups

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Definition

A **matrix group** over a field K is a subgroup of $GL_d(K)$.

- Matrix groups arise naturally in mathematics (and elsewhere).
- They are one of the basic ways of representing groups on a computer. In practice, they are usually given by finite sequences of generating matrices.

Computing with matrix groups

Computing with matrix groups over finite fields

- The “*Matrix Group Recognition Project*” has essentially been completed (≈ 20 years).
- Many problems have efficient solutions (up to basic oracles).
- Implementations exist (or are being developed) in GAP or MAGMA.

Computing with matrix groups over infinite fields

- Fundamental problems are undecidable (e.g. membership).
- Many computational problems are open.

Recent project (Detinko & Flannery, Eick, O’Brien, ...)

Develop and implement practical algorithms for basic computational problems related to finitely generated matrix groups over infinite fields.

Fundamental problems

Let $G \leq \mathrm{GL}_d(K)$ be given by finitely many generators.

Group-theoretic problems

- Decide if G is finite, soluble, nilpotent,
- Compute centralisers, normalisers,
- Find a finite presentation for G (if possible).
- ...

Geometric problems

- Decide irreducibility, primitivity, ... of G .
- Orbit-stabiliser problem.
- Decide membership in G of a given $x \in \mathrm{GL}_d(K)$.
- ...

Irreducibility testing

Definition

Let $G \leq \mathrm{GL}_d(K)$. Then G is **reducible** if there exists a proper G -invariant subspace of K^d . Otherwise, G is **irreducible**.

Computational tasks

- 1 Decide if G is irreducible.
- 2 If G is reducible, construct an invariant subspace.

Relevance

- Irreducibility is the most fundamental module-theoretic property.
- Reduction to irreducible groups is a common technique in the theory of matrix groups.
- Starting point of MGRP over finite fields.

Irreducibility testing

Computational tasks (again)

- 1 Decide if G is irreducible.
- 2 If G is reducible, construct an invariant subspace.

The state of the art

- Irreducibility of matrix groups over finite fields can be tested using the MEAT-AXE algorithm. (Parker 1984, Holt & Rees 1994, ...)
- Irreducibility of finite matrix groups over the rationals can be tested effectively (Nebe & Steel 2009). Implemented in MAGMA.

Computing with nilpotent matrix groups

- Nilpotent matrix groups have been studied extensively (Suprunenko).
- They have been shown to be well-suited for computations.
- Nilpotency and finiteness can be tested (Detinko & Flannery 2008).
- In (Detinko & Flannery 2006), an algorithm which simultaneously tests irreducibility and primitivity of nilpotent matrix groups over finite fields was developed.

This talk

Let K be a number field. We describe an algorithm for deciding irreducibility of f.g. nilpotent matrix groups over K . In the case of finite nilpotent groups, we obtain a fully constructive algorithm. Time permitting, we also consider primitivity testing of finite nilpotent groups.

Matrix groups and algebras

Definition

Let $G \leq \mathrm{GL}_d(K)$.

- The **enveloping algebra** $K[G]$ of G is the subalgebra of $M_d(K)$ generated by G .
- If $K[G]$ is semisimple, then G is **completely reducible**.
- If $K[G]$ is simple, then G is **homogeneous**.

Fact

G irreducible $\Rightarrow G$ homogeneous $\Rightarrow G$ completely reducible.

Fact (Detinko & Flannery 2008)

If G is f.g. nilpotent, then we can either prove that G is completely reducible or we can construct a proper $K[G]$ -submodule.

The strategy

Let $G \leqslant \mathrm{GL}_d(K)$ be f.g., nilpotent, and completely reducible.

Goal: decide irreducibility of G .

We proceed as follows:

- 0 The case that G is abelian is easily treated.
- 1 Find an abelian normal subgroup $A \triangleleft G$ which is either inhomogeneous or homogeneous and maximal (i.e. $A = C_G(A)$).
- 2 If A is inhomogeneous, then we can either prove reducibility of G or we reduce irreducibility testing to a problem in smaller dimension.
- 3 If A is homogeneous and maximal, then we can use computational Galois cohomology to decide irreducibility of G .

Step 1: constructing abelian normal subgroups

Let $G \leq \mathrm{GL}_d(K)$ be finitely generated, nilpotent and completely reducible.

Using congruence homomorphisms

We can find a homomorphism $G \xrightarrow{\pi} H$ onto a finite group H with the following property: a subgroup $A \leq G$ is abelian iff A^π is abelian.

Explicitly:

- Let R be the ring of integers of K .
- Choose an odd unramified prime $\mathfrak{p} \nmid R$ such that $G \leq \mathrm{GL}_d(R_{\mathfrak{p}})$.
- Take π to be the natural map $G \rightarrow G \bmod \mathfrak{p} \leq \mathrm{GL}_d(R/\mathfrak{p})$.
- By a theorem of Suprunenko and basic ANT, $\mathrm{Ker}(\pi)$ is torsion-free.
- Another result of Suprunenko implies that $[G, G]$ is finite.
- Hence, if A^π is abelian, then $[A, A] \leq \mathrm{Ker}(\pi) \cap [G, G] = 1$.

Step 1: constructing abelian normal subgroups

Goal: find an abelian $A \triangleleft G$ which is inhomgs or homgs and maximal.

Fact

A completely reducible abelian $A \leq \mathrm{GL}_d(K)$ is homg iff $K[A]$ is a field.

Fact (Dixon; Eberly)

For a completely reducible abelian $A \leq \mathrm{GL}_d(K)$, “most” elements $x \in K[A]$ satisfy $K[A] = K[x]$.

Let $G \xrightarrow{\pi} H$ be as on the previous slide. Using a presentation of H we find generators of $\mathrm{Ker}(\pi)$; note that $\mathrm{Ker}(\pi) \leq \mathrm{Z}(G)$.

- ① Let $B \triangleleft H$ be abelian.
- ② If $B^{\pi^{-1}}$ is inhomogeneous, then stop.
- ③ If B is maximal abelian, then stop.
- ④ Enlarge $B < C_H(B)$ and go to ②.

Step 2: reduction

Goal: make use of an inhomogeneous normal subgroup.

Theorem (Clifford 1937 + Detinko & Flannery 2006)

Let $G \leq \mathrm{GL}_d(K)$ be completely reducible, $N \triangleleft G$, and $K^d = U_1 \oplus \cdots \oplus U_r$ be the homgs decompn over $K[N]$. Then G is irreducible if and only if

- ① G acts transitively on $\mathcal{U} = \{U_1, \dots, U_r\}$, and
 - ② $\mathrm{Stab}_G(U_1)$ acts irreducibly on U_1 .
- \mathcal{U} can be easily computed if N is abelian.
 - We may test transitivity of G on \mathcal{U} using perm. group algorithms.
 - The action of $\mathrm{Stab}_G(U_1)$ on U_1 can be computed using linear algebra.

Reduction

If N is inhomogeneous (i.e. $|\mathcal{U}| > 1$), then we can either prove reducibility of G or we continue irreducibility testing in smaller dimension.

A reminder: crossed products

Let L/Z be a finite Galois extension of number fields, $\Gamma = \text{Gal}(L/Z)$, and $\phi \in Z^2(\Gamma, L^\times)$. Define

$$L \star_\phi \Gamma = \bigoplus_{\sigma \in \Gamma} u_\sigma L$$

with multiplication $au_\sigma = u_\sigma a^\sigma$ ($a \in L, \sigma \in \Gamma$) and $u_\sigma u_\tau = u_{\sigma\tau} \cdot (\sigma, \tau)\phi$. The algebra $L \star_\phi \Gamma$ is the **crossed product** of L by Γ determined by ϕ .

Facts

- $\mathcal{A} = L \star_\phi \Gamma$ is a central simple Z -algebra.
- $\text{index}(\mathcal{A}) = L$ -dimension of the irreducible \mathcal{A} -module.
- $\text{index}(\mathcal{A}) = \text{order of } [\phi] \in H^2(\Gamma, L^\times)$. (Brauer-Hasse-Noether 1932)
- The order of $[\phi]$ can be determined algorithmically. (Fieker 2009)

Step 3: cohomology

Goal: decide irreducibility of G if $A \triangleleft G$ is homogeneous and max. abelian.

Proposition

Let $G \leq \mathrm{GL}_d(K)$ be nilpotent and let $A \triangleleft G$ be max. abelian and homgs.

- 1 Let $L = K[A]$. Then G/A acts faithfully on L by conjugation.
- 2 Let $Z = L^{G/A}$. Then $K[G] \cong_Z L \star G/A$ in the natural way.

We may thus decide irreducibility of G as follows:

- 1 Construct $\phi \in Z^2(G/A, A)$ corr. to $1 \rightarrow A \rightarrow G \rightarrow G/A \rightarrow 1$.
- 2 Compute $m = \text{order of } [\phi] \in H^2(G/A, L^\times)$.
- 3 Check if $d = m|L : K|$.

Summary (deciding irreducibility)

Deciding irreducibility of a f.g. nilpotent group $G \leq \mathrm{GL}_d(K)$ (main steps):

- 1 Construct an abelian normal subgroup $A \triangleleft G$ which is inhomogeneous or maximal abelian and homogeneous.
- 2 If A is inhomogeneous, then either prove reducibility or reduce to smaller dimension and start again.
- 3 If A is maximal abelian and homogeneous, then compute the index of $K[G]$ and read off irreducibility of G .

Remark

We don't (in general) obtain a submodule in case 3.

Example: cyclic algebras

Suppose that K contains a primitive m th root of unity ζ_m . Let $\lambda, \nu \in K^\times$ and suppose that $X^m - \nu$ is irreducible over K . Let $\beta = \sqrt[m]{\nu}$ and define

$$G = \left\langle \underbrace{\begin{bmatrix} \beta & & & \\ & \beta \cdot \zeta_m & & \\ & & \ddots & \\ & & & \beta \cdot \zeta_m^{m-1} \end{bmatrix}}_u, \begin{bmatrix} 0 & 1 & & \\ & \ddots & \ddots & \\ & & 0 & 1 \\ \lambda & & & \end{bmatrix} \right\rangle.$$

Note that G has class 2. Regard G as a matrix group over K of degree m^2 . Then $A = \langle \lambda, u, \zeta_m \rangle$ is a homomorphism maximal abelian normal subgroup of G with $G/A \cong C_m$. It is well-known that $K[G]$ is a **cyclic algebra**.

Fact

Let m be a prime. Then G is reducible iff $N_{K(\beta)/K}(x) = \lambda$ has a soln.

The finite case

- We obtain a fully constructive algorithm for irreducibility testing of finite nilpotent matrix groups.
- We can handle a considerably larger class of fields of char. zero.
 - ▶ In theory: any field of characteristic zero with algorithms for polynomial factorisation and for solving $x^2 + y^2 = -1$.
 - ▶ In practice: number fields and rational function fields over these.
- The algorithm is quite practical. An implementation is available in MAGMA V2.17 and as a stand-alone package.
- Primitivity can be tested too.

The strategy (finite case)

Fact

If $A \leq \mathrm{GL}_d(K)$ is finite non-cyclic abelian, then A is inhomogeneous.

Strategy (based on (Detinko & Flannery 2006) for finite fields)

Let $G \leq \mathrm{GL}_d(K)$ be finite nilpotent.

- 1 Find a non-cyclic abelian normal subgp of G or prove that none exists.
- 2 In the first case, reduce.
- 3 In the second case, test irreducibility of G directly.

Theorem (Roquette 1958, ...)

Let G be finite nilpotent. All abelian normal subgroups of G are cyclic iff

- G_2 is cyclic or isomorphic to Q_8 or to D_{2^k} , SD_{2^k} , Q_{2^k} for $k \geq 4$, and
- $G_{2'}$ is cyclic.

Finding non-cyclic abelian normal subgroups

We can find $A \triangleleft G$ which is non-cyclic or cyclic and max. abelian.

Lemma

Let G be a finite nilpotent group such that $[G, G]$ is cyclic. Define $H = C_G([G, G])$.

- 1 If $H_{2'}$ is cyclic and H_2 is cyclic or $H_2 \cong Q_8$, then all abelian normal subgroup of G are cyclic.
- 2 Suppose that $H_p \not\cong Q_8$ is non-abelian. Then $\langle Z(H_p), h \rangle$ is a non-cyclic abelian normal subgroup of G for some $h \in H_p$.

Proof.

- 1 Follows from (Berger, Kovács, Newman 1980).
- 2 Note that $\text{class}(H_p) = 2$. If $\langle Z(H_p), h \rangle$ were cyclic for all $h \in H_p$, then H_p would contain a unique subgroup of order p . But then H_p would be cyclic or generalised quaternion, which is impossible. \blacklozenge

Enter $x^2 + y^2 = -1$

Let all abelian normal subgps of a non-abelian $G \leq \mathrm{GL}_d(K)$ be cyclic.

- Let $A \triangleleft G$ be cyclic of index 2. We may assume that A is homgs.
- Let $Z = Z(K[G]) = K[A]^G$. We find that

$$K[G] \cong_Z \left(\frac{-1, \pm 1}{Z} \right) = Z(\sqrt{-1}) \star C_2.$$

Lemma

- ① If G_2 is (semi)dihedral, then G is irreducible iff $d = 2|Z : K|$.
- ② Let G_2 be quaternion. If $x^2 + y^2 = -1$ is soluble in Z , then G is irreducible iff $d = 2|Z : K|$. Otherwise, G is irreducible iff $d = 4|Z : K|$.

Enter $x^2 + y^2 = -1$

Example

We have $K[G] \cong \left(\frac{-1, -1}{K} \right)$, where

$$G = \left\langle \begin{bmatrix} \cdot & 1 & \cdot & \cdot \\ -1 & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & -1 \\ \cdot & \cdot & 1 & \cdot \end{bmatrix}, \begin{bmatrix} \cdot & \cdot & 1 & \cdot \\ \cdot & \cdot & \cdot & 1 \\ -1 & \cdot & \cdot & \cdot \\ \cdot & -1 & \cdot & \cdot \end{bmatrix} \right\rangle \leq \mathrm{GL}_4(K).$$

Hence, G is reducible iff $x^2 + y^2 = -1$ is soluble in K . If this is the case, then finding a proper $K[G]$ -submodule is equivalent to finding a solution of $x^2 + y^2 = -1$.

Primitivity testing

Definition

Let $G \leq \mathrm{GL}_d(K)$ be irreducible. If there exists a non-trivial decomposition

$$K^d = U_1 \oplus \cdots \oplus U_r$$

permuted by G , then G is **imprimitive**. Otherwise, G is **primitive**. A subspace U_i is a **block** and a subgp $\mathrm{Stab}_G(U_i)$ is **block stabiliser** for G .

Computational tasks

- 1 Decide if G is primitive.
- 2 If G is imprimitive, construct a system of imprimitivity.

Elementary facts

Facts

Let $G \leq \mathrm{GL}_d(K)$ be irreducible.

- $H < G$ is a block stabiliser iff there is an irreducible $K[H]$ -submodule $U < K^d$ with $d = |G : H||U : K|$.
- G is imprimitive iff some max. subgp of G is a block stabiliser.
- If $H < G$ has index 2, then H is a block stabiliser iff H is reducible.
- If G is primitive, then all normal subgroups of G are homogeneous.

Maximal subgroups

Let $G \leq \mathrm{GL}_d(K)$ be finite, nilpotent, and irreducible.

- If $A \triangleleft G$ is non-cyclic abelian G , then G is imprimitive.
- We obtain a reduction to the case that all abelian $A \triangleleft G$ are cyclic.
- The abelian case is easily treated (again).
- Let $A \triangleleft G$ be cyclic of index 2. We may assume that A is irreducible.
- It suffices to test if any of the max. subgps of G is a block stabiliser.
- The interesting ones correspond to the prime divisors of $|G|$.

Lemma

$|G| = \mathcal{O}(d^{1+\varepsilon})$ for $\varepsilon > 0$.

Maximal subgroups

Proposition

- ① *Let $H < G$ have prime index p (+ conditions for $p = 2$, e.g. $A \neq H$). Suppose that one of the following conditions is satisfied:*

- ▶ G_2 is (semi)dihedral,
- ▶ $|G_2| \geq 32$, or
- ▶ p is odd.

Then H is a block stabiliser iff $|K[A] : K[A^p]| = p$.

- ② *If $G_2 \cong Q_8$, then subgroups of index 2 of G are irreducible.*

- ③ *Suppose that $Q_8 \times C_m \cong H < G \cong Q_{16} \times C_m$ for odd m .*

Then H is a block stabiliser iff $|K[A] : K[A^2]| = 2$ and the following condition is satisfied: $\text{ord}(2 \bmod m) | K_p : \mathbb{Q}_2 |$ is even for all $p \mid 2$.

We can thus test primitivity of G by looping over its maximal subgroups and testing the conditions in the proposition.